SBL: A Simple Programming Language for Business

Craig E. Ward, Paul Bull CMSI 585 Programming Languages Loyola Marymount University

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1 Introduction

SBL, Simple Business Language, is intended as a straight forward, simple, imperative language for programming business applications. The features of SBL are those that the designers feel are most appropriate for business applications. These include:

- Basic mathematical operations
- Strings as a native type
- Concurrent processing
- Exception handling
- A library of typical accounting functions
- A library of standardized database connectivity routines

SBL should provide programmers with the materials needed to model the steps of a business processes. Such processes are usually, in-and-of themselves, simple and straight forward. The complexities arise from the way each of these steps interact. By allowing a focus on the steps instead of the overall process, SBL should serve the business community better than a language that allows a technical team to attack enterprise issues.

All of the features of SBL can be found in other languages. SBL has been designed from an *a la carte* survey of many programming languages, especially C/C++, Java, Perl, and the much maligned Pascal. The goal is to design a language that allows for a variety of programming styles without allowing so much variety that it becomes easy for the code to be confusing. The *References* section lists the languages surveyed.

2 Syntax and Semantics Overview

2.1 Microsyntax

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The source code for a SBL program is a sequence of Unicode characters. Comments start with "//" and extend to the end of line. Language tokens are comprised of the longest possible match of characters. Tokens are separated by white space characters or comments.

Identifiers are single token strings of alphabetic characters, digits, or the underscore character starting with an alphabetic character that are not reserved words. SBL has the following reserved words:

boolean catch char do else false float int if null record secret string synchronized task true try until use void while

Integer literals are a sequence of digits. Floating point literals are identified as a sequence of digits followed by a period (".") followed by at least one digit. Scientific notation for floating point numbers is not part of SBL.¹

String literals are a sequence of printable characters enclosed between double quotes.

The following sequences can be used to encode additional, special characters:

The language also lacks bit operators; SBL should not be used for system programming.

2.2 Macrosyntax

Below is the EBNF description of SBL syntax. Brackets indicate an optional item;

tokens in curly braces appear zero or more times. Tokens in lowercase are reserved

words.

```
PROGRAM ::== MODULE
MODULE  ::= \{ use ID 'i' \} \{ FUNCTION\}FUNCTION ::== [secret] PROFILE BODY
BODY ::== '{' [TASKBLK] SERIALBLK '}'
PROFILE ::== (TYPE | void) ID '(' PARAMS ')'
TYPE ::== (boolean|int|float|char|string|RECORD|ID) { '[' ']' }
RECORD ::== record '{' TYPE ID ';' {',' TYPE ID ';' } '}'
PARAMS  ::= = [TYPE [ '&'] ID { , TYPE [ '&'] ID }TASKBLK ::== {task '{' FUNCTION '}'}
SERIALBLK ::== { STATEMENT | '{' SERIALBLK '}' } | try '{' SERIALBLK
'}' catch '{' SERIALBLK '}'
STATEMENT ::== TYPE ID ':=' EXP '; '
      | ID ':=' EXP ';'
      | EXP ';'
      | if '(' EXP ') '{' SERIALBLK '}' [else if '{' SERIALBLK '}'] [
else '{' SERIALBLK '}']
      | return [EXP] ';'
      | while '(' EXP ')' '{' SERIALBLK '}'
      | do '{' SERIALBLK '}' until '(' EXP ')' ';'
EXP ::== EXPO \{ | \ | \ EXPO \}EXP0 ::== EXP1 { '&&' EXP1 }
EXP1 ::== EXP2 \{ RELOP EXP2 \}EXP2 ::== EXP3 \{ ADDOP EXP3 \}EXP3 ::== EXP4 { MULOP EXP4 }EXP4  ::= [PREFIXOP] EXP5
EXP5 ::== LITERAL | ID | ID '(' ARGS ')' | '(' EXP ')'
LITERAL ::== null|true|false|INTLIT|FLOATLIT|CHARLIT|STRINGLIT
ARGS ::== [ EXP \{ ' , ' EXP \} ]
ADDOP ::== ' +' | ' -'
MULOP ::== '*' | '/' | '%'
PREFIXOP ::== '-' | '!'
```
2.3 Semantics

An SBL program is a set of one or more functions. Below is an outline of the semantics of the language.

2.3.1 Abstract Syntax

The abstract syntax tree makes our defining of semantics easier by allowing us to use abbreviated aliases for terms that we defined in our syntax. This makes describing the semantic functions of the language more compact and easier to read.

```
op: PREFIXOP|ADDOP|MULOP
lit: LITERAL -> null | true | false | INTLIT | FLOATLIT | CHARLIT |
STRINGLIT
id: ID
e: EXP -> e'||'e | e && e | e+e| e-e| e*e | e/e | e%e
s: STMT
sb: SERIALBLK
tb: TASKBLK
parms: PARAMS
type: TYPE
bod: BODY
pro: PROFILE
fun: FUNCTION
mod: MODULE
prog:PROGRAM
```
2.3.2 Semantic Domains

The domains of a program consist of modules which carry out behavior that effects the state of a program, files that may be written to or read from and the store which is a mapping of variables to integer values.

```
_: module = int list
_: file = int list
\_: store = ID \_ int
\therefore state = store x {modules} x {files}
```
2.3.3 Semantic Functions

A program is a modification from state to state, state to file, file to state or file to file. As mentioned before a program is comprised of one or more modules. All computation starts and ends with the *begin* function which executes statements and calls other functions. The execution of a statement or sequence of statements changes the state of

computation. The execution of expression or assignments changes the store which is a

mapping of IDs to integer values.

```
P: PROGRAM _ MODULE _ {MODULE}
M: MODULE _ state _ state _ {files}
F: FUNCTION _ state _ state
S: STATEMENT _ store _ store
E: EXP _ store _ LITERAL
SB:SERIALBLOCK _ store _ store
P = | = int(M \mod, {M \mod})M = (fun) +F (e1, r, ,en) = E e (S s, F fun (x.undef))SB(s1.....sn) = S s1 = ...S sn =S( id := e ) = [E(e) | id ]S(e) = int[E(e)]S(if e sb1 else sb2) = SB sb1 _ | SB sb2S(if e sb1 elsif sb2 else sb3) = S s1 - | S s2 - | S s3 -S(\text{return } e) = E(e)S(while e sb) _ = E e _ _ S(while e sb)(Sb sb _) | _
S(do sb while e) _ = SB sb _ E e _(S(do sb while e) | _)
E(\text{true}) = \text{lit}E(false) = litE(e1<e2) = E e1 = < E e2 =
E(e1<=e2) = E e1 = \leq E e2 =E(e1 > e2) = E e1 = > E e2 =E(e1>=e2) = E e1 _ >= E e2 _
E(e1=e2) = E e1 = = E e2 =
E(e1!= e2) = E e1 = \neq E e2 =
E(e1 \& k \cdot e2) = E e1 = E e1E(e1 | e2) = E e1 - V E e2E(-e) = -(E e)E(e1 + e2) = E e1 = + E e2 =
E(e1 - e2) = E e1 - E e2E(e1 * e2) = E e1 - E Ee2E(e1 / e2) = E e1 / E e2 -E(e1 % e2) _ = E e1 _ % E e2 _
```
3 Names, Declarations, and Scopes

The initial function is called *begin*, which takes two arguments from the invoking environment: an array of string arguments and an integer count of arguments in the array.

The profile of this function is:

```
int begin(string [] arguments, int count)
```
The names used for the formal parameters here are examples only. An SBL program may use any names; it is the types that are required.

Multiple functions may be declared in a file. Unless specifically indicated by the *secret* reserved word, functions are globally visible. No relation exists between the name of the object, *e.g.* file, containing the function and the function itself.

Objects may be declared at any point in a function body, but may not be used until

after declared. ² Objects have lexical scope from the point of declaration to the end of the

compassing SERIALBLK. A name may not be reused in a nested SERIALBLK

construct.

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4 Types

SBL provides a set of types most useful to business data processing as well as some

generally useful for programming.

- The type *int* for integers, the range being implementation dependent.
- The type *float* for real, floating point, numbers. An implementation only needs to be able to support a range sufficient to represent $1/1000th$ of a unit value.
- The type *string* representing alphanumeric sequences
• A *record* is a named collection of other types.
- A *record* is a named collection of other types.
- An *array* is a named sequence of a single type.
- A *boolean* may only have a value of *true* or *false*.
- The *char* type may hold a single Unicode character.
- Functions are the base type of SBL programs.

All types are first class types. An object of any type may be passed as an argument to

any function and returned from any function. By default, the types *int*, *float*, and *boolean*

are passed by value. The types *string*, *record*, *array*, and *function* are passed by

 2^{2} In SBL, object is used to refer to anything that can be named in the language, not as the term is used in object oriented systems.

reference. Pass by value may be overridden by use of the "&" character in the list of formal parameters to a function.

5 Expressions and Statements

Statements produce side effects. An expression produces a value for one of the SBL types. Expressions in actual parameters are evaluated using lazy evaluation. An expression in a simple statement will short circuit for boolean values.

5.1 Examples

• int index := 0;

The identifier *index* is declared and then initialized with the integer value 0.

• string item_description;

The identifier *item_description* is declared. The value of the object is undefined.

• overdrawn $:=$ true:

A previously declared boolean, *overdrawn*, is assigned the value *true*;

• if (count ≤ 100) { order_items(product); }

The integer value of *count* is compared to 100; if the expression evaluates to *true*, call the function *order_items* with the argument *product*, otherwise do nothing.

- while (next_account != null) {overdrawn = check_acct_status(next_account);
	- }

Until the value of *next_account* becomes null, call function *check_acct_status* with *next_account* as an argument. Assign the returned value to *overdrawn*. The type of next_account must match the expected type of the *check_acct_status* formal parameter and the function's return type must match the type of *overdrawn*.

6 Subroutines and Modules

6.1 Functions

A function is defined by supplying the keyword *secret* if the function is only to remain accessible within its own file, if the *secret* keyword is omitted, the function is accessible to other files that may be linked to the functions module. Next the return type is specified for the return value of the function or *void* if the function does not return a value. Return types include *boolean, float, int, char, string, record*, or a user-defined type. Next the name of the function is supplied. After that, the parameters of the function are listed within parenthesis. Parameters are specified by type and then name followed by a comma if an additional parameter follows. Lastly the body of the function is supplied within curly braces. If a return type was specified for the function a return statement returning the appropriate type must be specified at the end of the function or at the end of each block of a sequence of conditionals where appropriate. The general skeleton for function definition is as follows:

```
["secret"] return type | "void" function-name '(' param1,…..paramN')'
'{'
 body
[return statement]
'}'
```
Examples are as follows:

1.)A function only accessible within its own module:

```
…
secret void modifyFields (int x, char y, float z)
{
 intfield = x;charfield = y;
 floatfield = z;
…
}
```
Where *intfield*, *charfield*, and *floatfield* would be variables of the function that called modifyFields or global variables of the module modifyFields is defined in.

2.) A function that is externally accessible and has a return type float addFloats(float x, float y) { return $x + y$; }

6.2 Modules

6.2.1 Module Creation

A module is defined by creating a file with a list of function definitions in it. The module file is begun with a label within tags that identifies the name of the module and optionally the directory in which it is to be stored in. The file is saved with the *.mod* suffix. The basic structure is as follows:

```
'<'[directoryname] ['<='] modname'>'
 …
     fun1( ){…}
     fun2( ){…}
 ….
     funN( ){…}
```
If the module name is specified with no directory name then it is assumed that it will be stored in the same directory as any file that might link to it. Otherwise, acceptable directory names would be, for example, *fio* for the directory of the text input output library or the names of any other libraries of the standard library. If a directory name is supplied that does not exist in the standard library a new directory with that name is created within the standard library.

Examples of module definition are as follows:

1.) An example of a module that is only accessible to programs within its own directory.

```
<mymodule>
    secret void fun1 ( int x, int y) \{...\}float fun2( float z, float v)\{...\}….
```

```
void funN( )\{...\}
```
2.) An example of a module that would be stored in the *fio* library

```
<fio <= txtmanip>
    string rotateTxt(string x){…}
    string clipTxt(string x, int begin, int end){...}
 …
```
3.) Lastly we have an example of a module that will be stored in a newly created subdirectory of the standard library.

```
<transactions <= transmod>
     void updateTransaction (record trans){…}
    void deleteTransaction(record trans){…}
   void modTransaction(record trans){…}
 …
```
transactions is not a library included within the standard library but would now be included with file **transmod.mod** containing the updateTransaction, deleteTransaction, modTransaction etc. functions available through the newly created *transactions* library.

6.2.2 Utilizing Modules

In order for a program to utilize functions of a particular module, the intent to use that module must be declared somewhere before the functions are defined in that program. The module is utilized by supplying the *use* keyword. If the module file is stored within one of the libraries of the standard library then the name of that library following *use* will suffice. If the module is a stand-alone module then the name of the module file in tags must follow *use.* Also, as mentioned earlier, in this case the module file and the file of the program using it must be in the same directory. Referring to modules in examples 1 and 2 of section 6.2.1 above, here are examples of how to use modules.

> 1.) An example using the mymodule.mod file used earlier: … use <mymodule.mod>; … …

```
begin( )
{….
   fun2(10.2,99.99);
}
```
Once more, it is important to remember that mymodule must be in the same directory as this program's file.

2.)The following is an example of a program that wants to call functions in txtmanip.mod file. … use \langle fio $>$; … begin() $\{\ldots$ rotateText("racecar"); … }

In this case since txtmanip.mod is in the *fio* library so only *fio* needed to follow *use.*

7 Concurrency

7.1 Threads

SBL adopts a style of concurrent processing similar to Ada.

A function declaration may include one or more TASKBLKs, each defining a thread of execution. The TASKBLK contains one or more functions. When the SERIALBLK, of the function begins to execute, the initial function of each TASKBLK also begins. The initial function of a TASKBLK may wait for arguments to be passed to it from the SERIALBLK. A TASKBLK function expecting arguments may be called from the SERIALBLK. Calling a TASKBLK function that does not expect arguments is a compiler error.

7.2 Synchronization

Objects shared between concurrently processing tasks can be enclosed in a *synchronized* block to coordinate access similar to the way Java handles the issue. For example:

```
synchronized { int global_counter; }
```
The SBL implementation must provide coordinated access to the variable *global_counter* using the available resources of the host environment.

8 Standard Library

SBL implementations provide a basic set of library calls. Library routines are included in a module by using the use statement. The use statement may appear at any point in a module. For example:

use <fio>; // include the standard file i/o routines

8.1 FIO: File Input-Output

The File Input-Output library provides routines for basic text and binary input and output. A record is defined that contains the data needed by the host implementation for controlling the location of reads or writes in a file. Some function must be called within a try/catch construct.

Include this line to use this library: use \leq fio>;

Functions defined in this library are:

8.2 NIO: Network Input/Output

The Network Input/Output library provides some basic routines for using network

communication resources. The details are implementation dependent. These routines

must be used within a try/catch block. NIO handles are compatible with FIO handles. FIO

handles can not be used with these routines.

Include this library with: use \langle nio>;

8.3 Currency

The currency library provides some basic math routines to aid in the processing of

dollars and cents. Include this library module with: use <currency>;

8.4 SDBC: SBL Database Connectivity

The SDBC library provides a standard set of functions for interacting with database

management systems. Include this library with: use \leq sdbc>;

These functions must be called within a try/catch block.

8.5 Utility

The utility library provides routines that may be of general use in an application. Use

this library with: use <utility>;

Appendix A

An example of code using synchronization and task blocks:

```
…
     synchronized { record account{
                            string firstname;
                            string lastname;
                            int customer_id;
                            float account_balance
                            }
                     }
     begin( ){ doTransactions( account);}
     void addtoAccount(record account,float amount)
     {
         account.balance += amount;
 }
     void printAccountBal(record account)
      {
          print("The balance to the account is now" +
account.balance);
     }
           void doTransactions( record account)
\{int i = 0;
              float amt = 2.0 while( i != 1000)
\{task{ addtoAccount(account, amt);}
           task{printAccountBal(account);}
 }
     }
```
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